Building an Efficient Key-Value Store in a Flexible Address Space

Chen Chen¹, Wenshao Zhong¹, Xingbo Wu²

¹ University of Illinois at Chicago
 ² Microsoft Research Cambridge

KV Stores

• The backbone of widely used Internet-based services

KV Stores

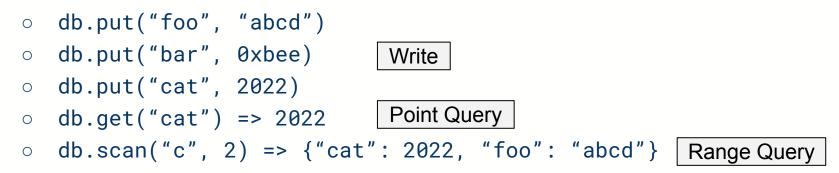
- The backbone of widely used Internet-based services
- Simple interface
 - db.put("foo", "abcd") Ο
 - db.put("bar", 0xbee) Ο
 - db.put("cat", 2022) 0
 - db.get("cat") => 2022 Ο
 - db.scan("c", 2) => {"cat": 2022, "foo": "abcd"} Range Query 0

Write

Point Query

KV Stores

- The backbone of widely used Internet-based services
- Simple interface



• Major KV stores manage **sorted data** for range queries

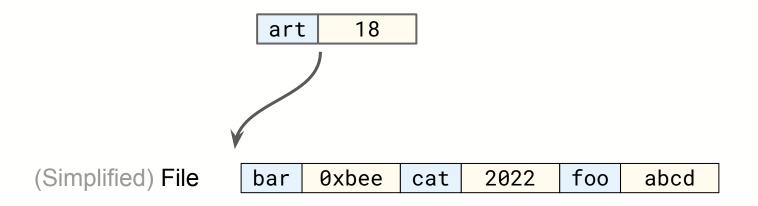
• Storing sorted data in structured files



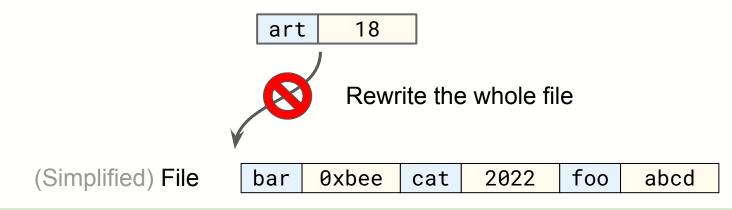
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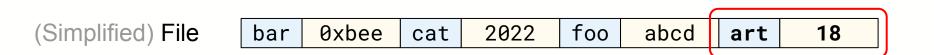
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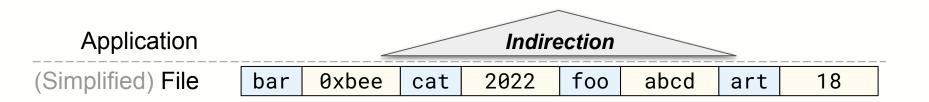
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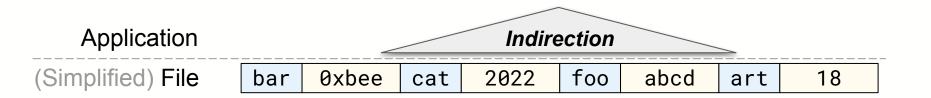


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 - Unordered file but sorted index for *all* keys

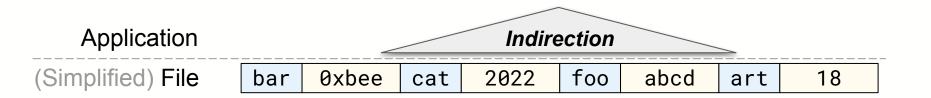
Application Indirection (Simplified) File foo 0xbee cat 2022 abcd 18 bar art

- Storing sorted data in structured files
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 - Unordered file but sorted index for all keys Maintaining large index

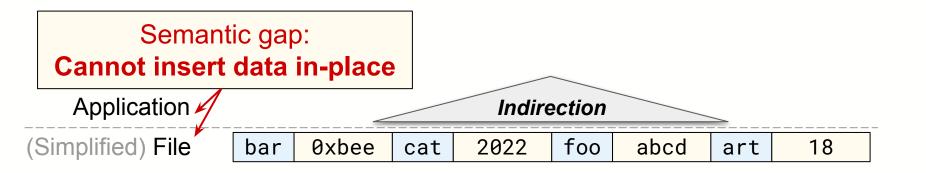
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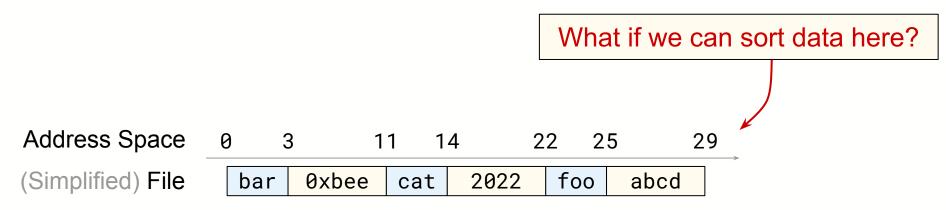


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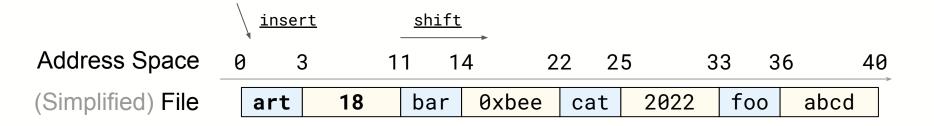


(::) Repeated *rewrites*

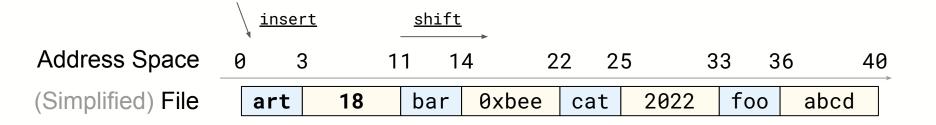
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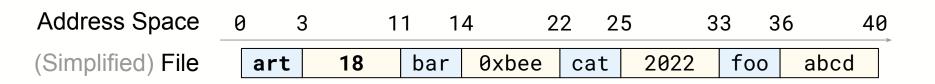
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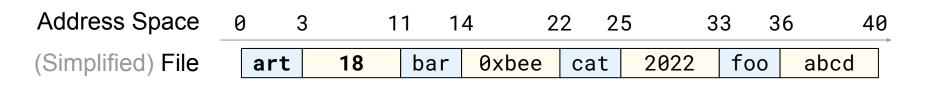
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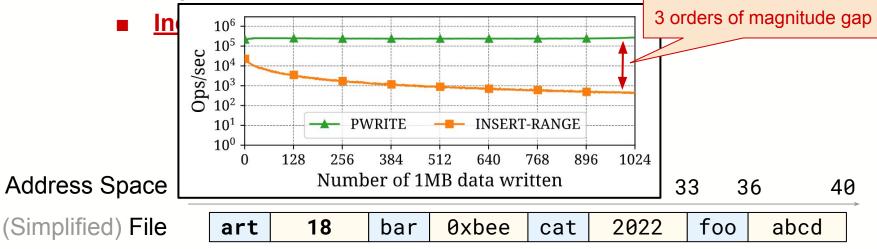


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Inefficient shifting in extent indexes

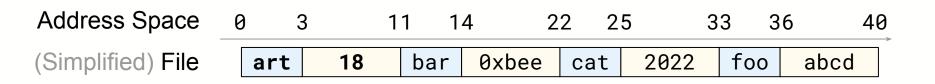
Updating O(N) extents' metadata Address Space 3 0 22 25 33 36 40 11 14 2022 foo (Simplified) File 18 0xbee cat abcd art bar

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Rigid block alignment requirements …



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 - Inability to track shifting data

The offsets can easily change

Address Space	0	3		11	14	4 2	2 2	5	33	36	40
(Simplified) File	ar	't	18	b	ar	Øxbee	cat	2022	fc	00	abcd

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A more flexible storage abstraction to manage sorted data!

Address Space	0	3	1	1 1	4 2	2 2	5 3	33 3	6 40	*
(Simplified) File	ar	't	18	bar	Øxbee	cat	2022	foo	abcd	

- Lightweight in-place insertions and deletions
 - Sorting data easily in the address space

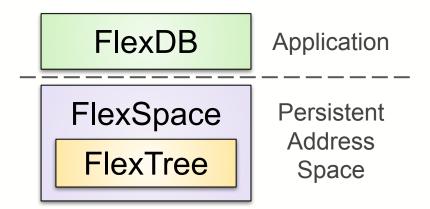
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- Lightweight in-place insertions and deletions
 - Sorting data easily in the address space
- Challenges => The Solution
 - Inefficient shifting in extent index => Index structure with efficient shifting
 - **Rigid block alignment requirements =>** No alignment requirements
 - Inability to track shifting data => Managing shifting data

- Lightweight in-place insertions and deletions
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- The solution:

Managing shifting data

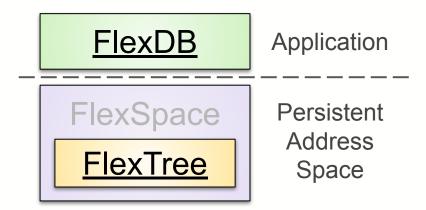
No alignment requirements Index with efficient shifting

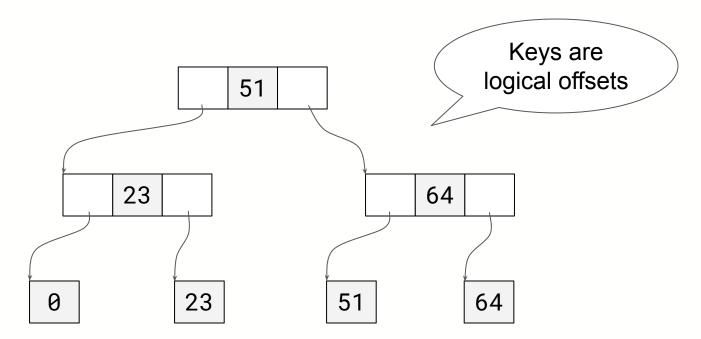


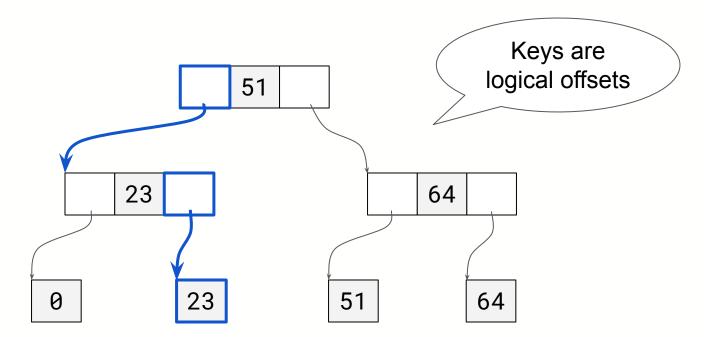
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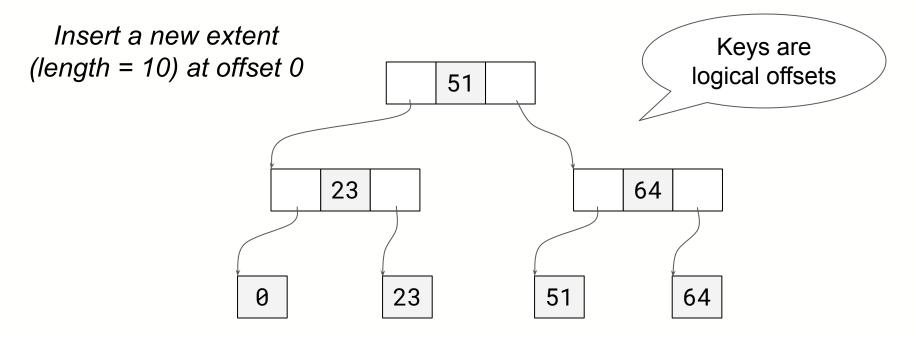
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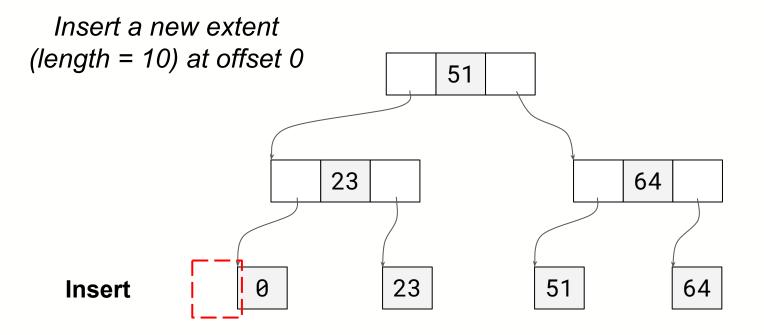
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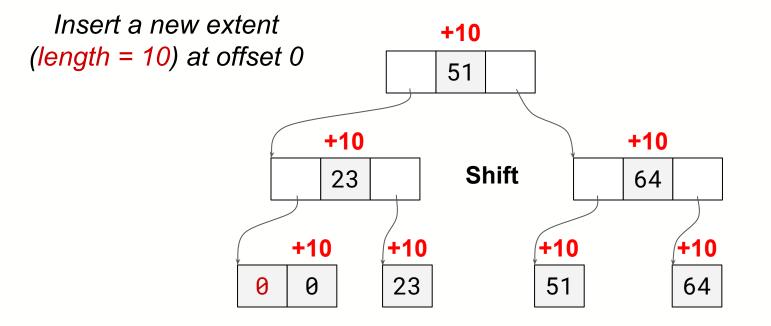






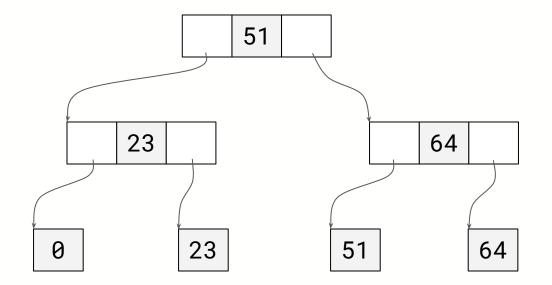






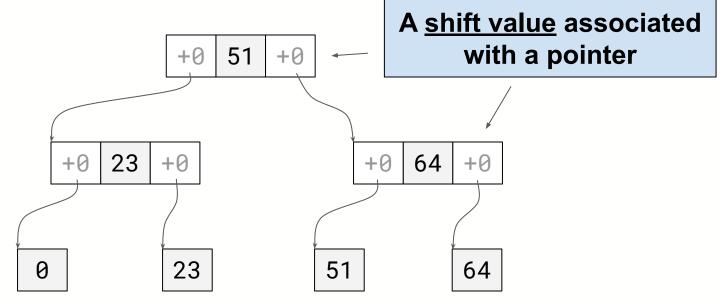
FlexTree: Structure

- An index structure derived from B⁺-Tree
 - A new metadata representation scheme



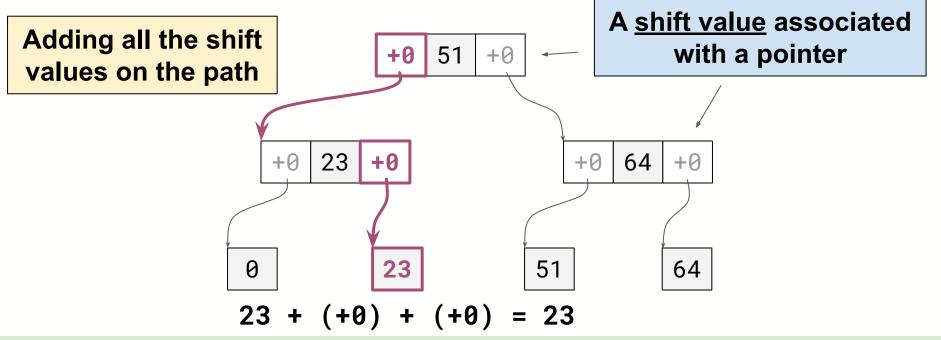
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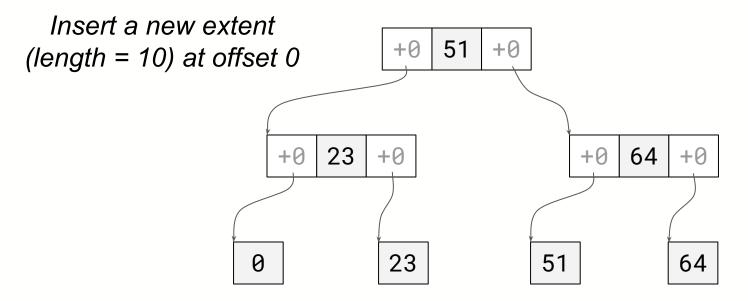


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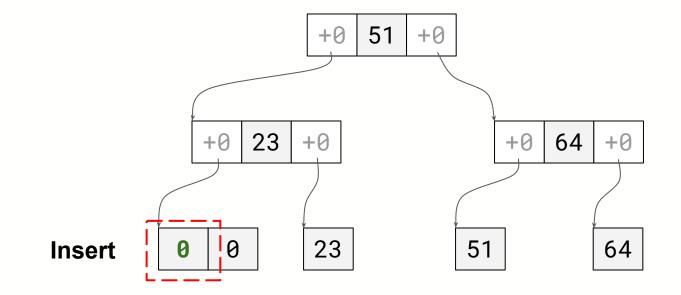
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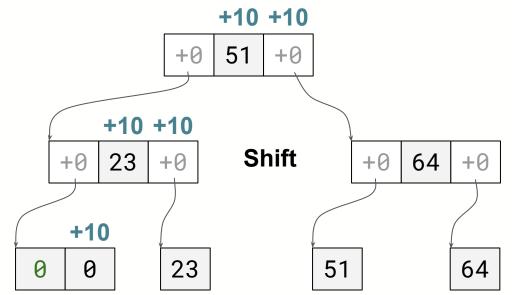
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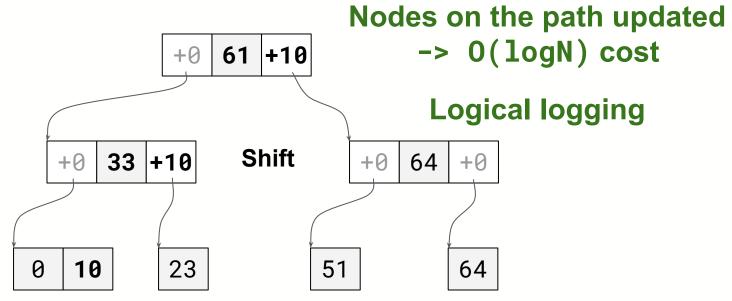
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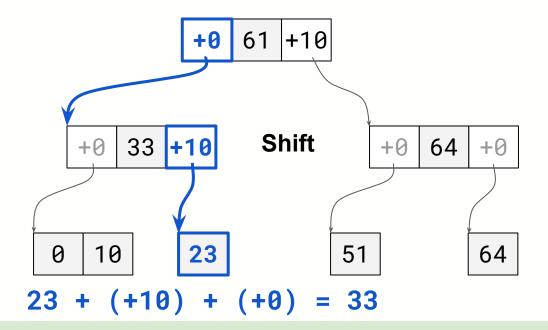
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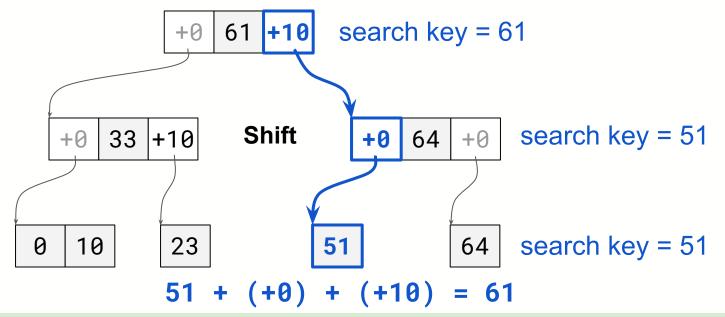
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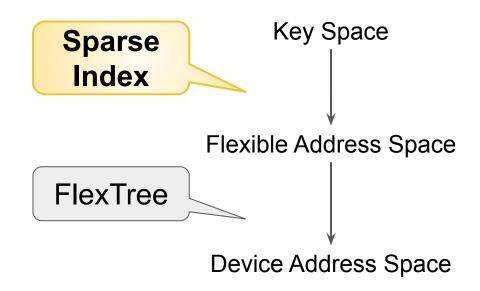


Based on FlexTree

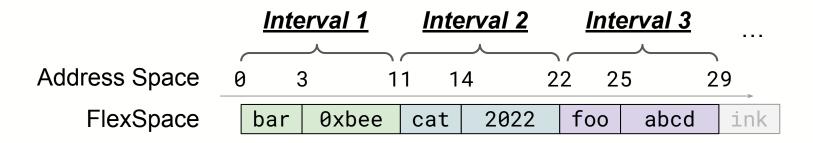
- FlexSpace: Log-structured data storage indexed by FlexTree
 - Supporting read/write/insert-range operations etc.
- FlexDB: Keeping all KV pairs sorted in a FlexSpace

FlexDB

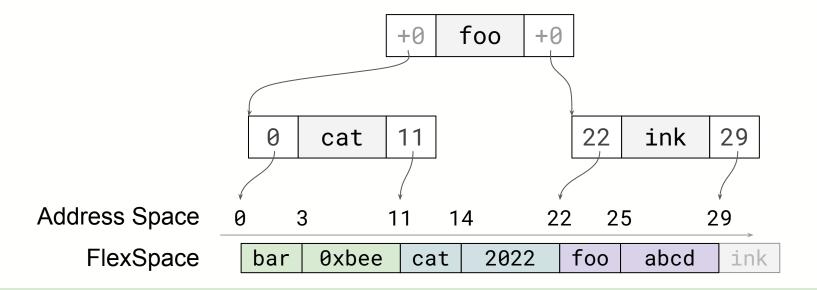
• Keeping all KV pairs sorted in a flexible address space



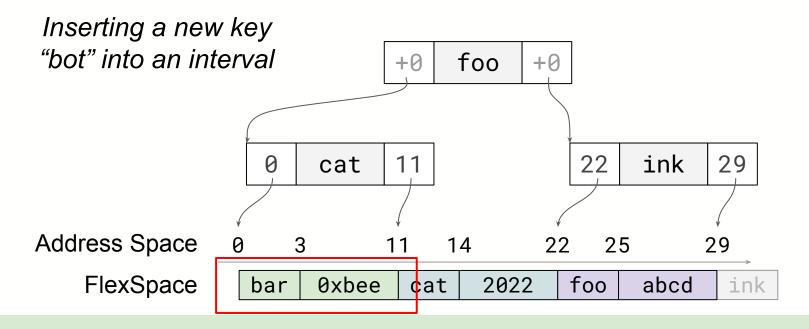
• Managing sorted keys in *intervals*



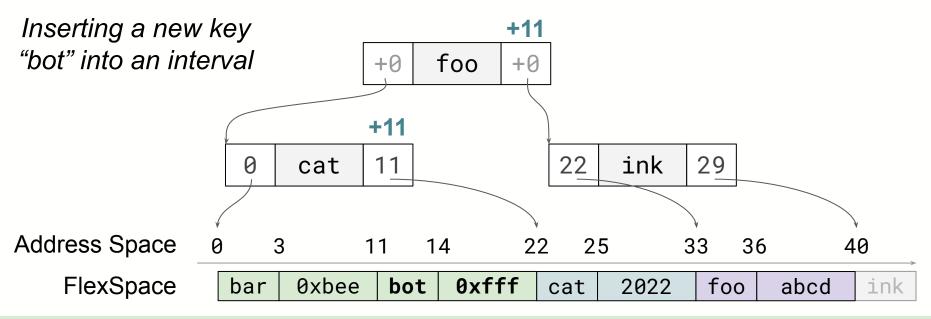
- Managing sorted keys in *intervals*
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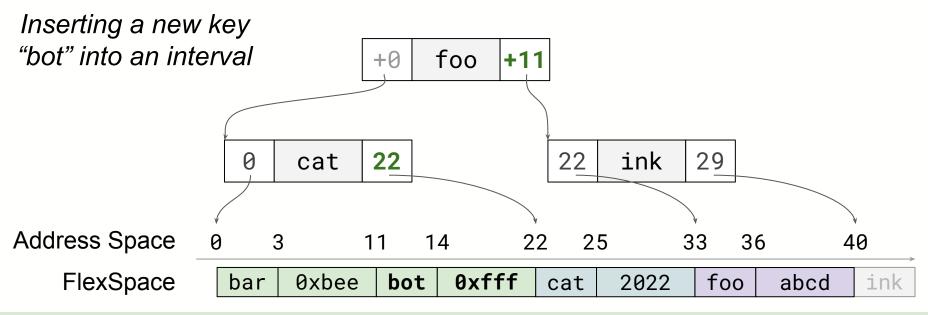
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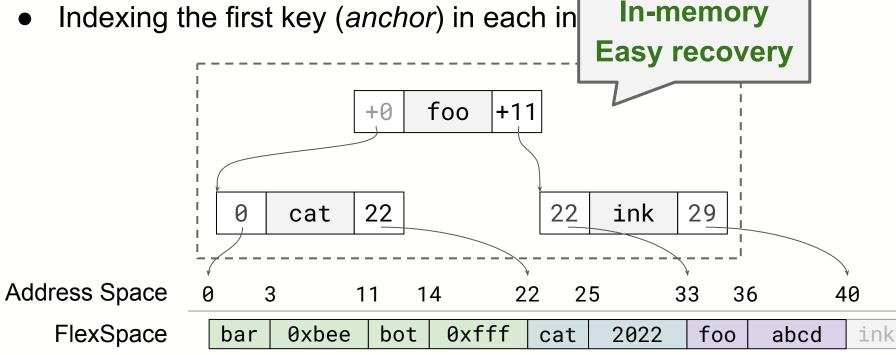


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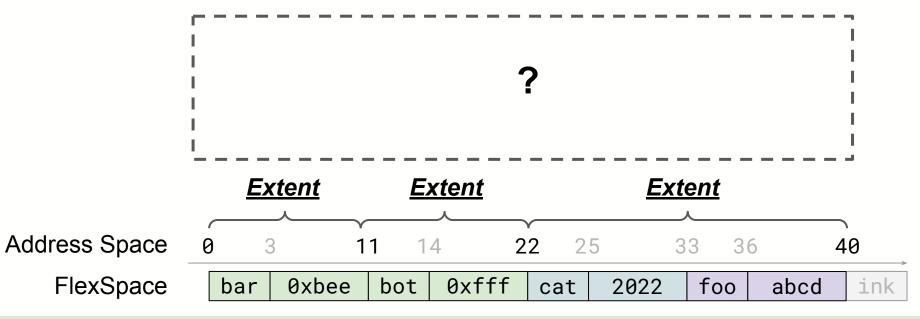
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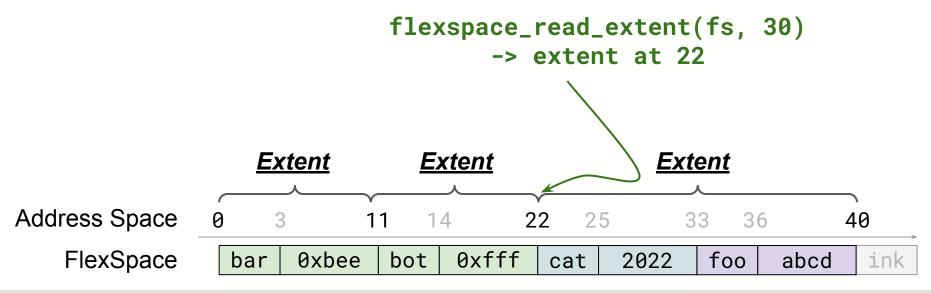
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Elastic

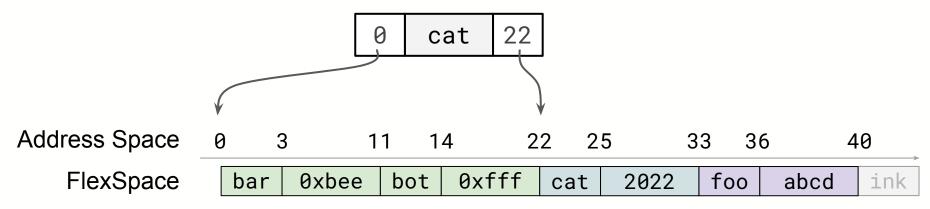
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- FlexTree enables lightweight data insertions in a flexible address space.
- FlexDB manages sorted data without using extra persistent indirections.

Evaluation: Setup

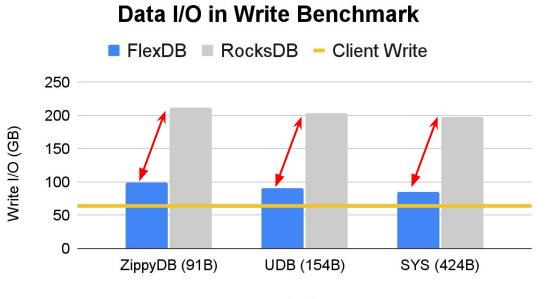
- Intel Xeon Silver 4210 w/ 10 cores
- 64GB RAM
- Optane 905P SSD
- Key-Value Sizes:
 - ZippyDB: 91 bytes *
 - UDB: 154 bytes *
 - SYS: 424 bytes **
- 4 Client Threads

* Zhichao Cao, Siying Dong, Sagar Vemuri, and David H. C. Du. "Characterizing, Modeling, and Benchmarking RocksDB Key-Value Workloads at Facebook". In: 18th USENIX Conference on File and Storage Technolo gies (FAST'20). 2020, pp. 209–223.

** Berk Atikoglu, Yuehai Xu, Eitan Frachtenberg, Song Jiang, and Mike Paleczny. "Workload Analysis of a Large-Scale Key-Value Store". In: SIGMETRICS Per form. Eval. Rev. 40.1 (2012), pp. 53–64.

Evaluation: Disk I/O

□ Write 64GB into an empty store; Zipfian distribution



KV Size

□ Starting from a 500GB UDB store; Zipfian distribution

RocksDB FlexDB 520K 857K 2.73M 552K 677K 169K 100 75 Throughput (%) 50 25 0 В С Ε F Α D 60 sec per Write Read Read Read Scan Read-modify-write workload (RMW) Mostly Mostly Only Latest Mostly

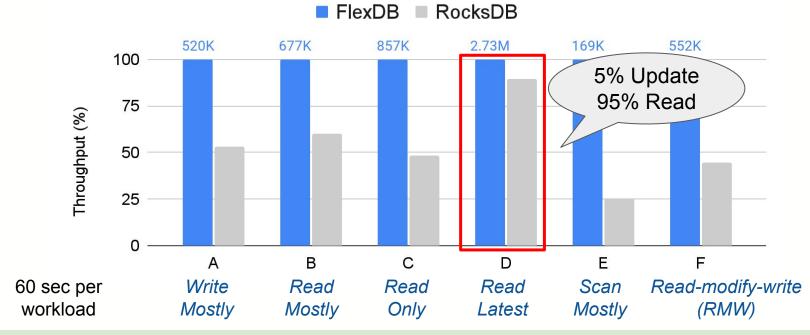
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Normalized YCSB Throughput

RocksDB

Summary



- Flexible address space enables lightweight data management.
- FlexDB achieves low WA and high throughput.

 The code of this project is available at: <u>https://github.com/flexible-address-space/flexspace</u>